

**IN THE CLAIMS:**

Please amend claims 1-13 and add claims 14-20 as follows:

## CLAIMS

1. (Currently Amended) A Mmethod of storing encrypted data in a random access memory, comprising the steps of:  
~~in which data words, which each comprise a predetermined number of data bits, are storable,~~  
~~characterized in that, before storage, an~~ \_\_\_\_\_ encryptiong of each data word by permutating each  
~~data bit of the data word using a permutation key to generate (M) is effected whereby a~~  
~~permuted data word; and~~  
~~\_\_\_\_\_ storing the permuted data word in the memory (P) with a predetermined number of data~~  
~~bits is generated from each data word (M), or from a data word (M) derived from this data word,~~  
~~by one to one permutation of the individual data bits (M[n-1] M[0]) using a first permutation~~  
~~key (Mp).~~

2. (Currently Amended) The Mmethod of according to Cclaim 1, in which where after the  
step of permutating, further comprising the step of substituting each ~~the individual data bits~~  
~~(M[n-1] M[0]) of the permuted data word (Mp) are substituted before storage using a first~~  
~~substitution key to generate a substitute data word, and where the step of storing comprises the~~  
~~step of storing the substitute data word in the memory in order to provide an encrypted data word~~  
~~(M').~~

3. (Currently Amended) The Mmethod of according to Cclaim 1, where the step of  
encrypting further includes the step of substituting each data bit of the unencrypted data word  
using a substitution key prior to the step of permutating to generate a substitute data word, and

where the step of permutating comprises the step of permutating each data bit of the substitute data word using the permutation key to generate the permuted data word~~in which before rearrangement the individual data bits of the data word (M) are substituted using a first substitution key (S) in order to provide a substituted data word.~~

4. (Currently Amended) The Mmethod according to one of the foregoing claims 1, where in which the permutation key (P) has a number includes a plurality of unique subkeys (P[n-1]-P[0]) corresponding to the number n of the data bits of the data word, and where each one of the subkeys includes a plurality of key bits, where the step of permutating each data bit in the data word using a permutation key further comprises the steps of:

\_\_\_\_\_ assigning each one of the subkeys to a corresponding one of the data bits of the permuted data word; and

\_\_\_\_\_ mapping each data bit of the unencrypted data word to a corresponding one of the data bits of the permuted data word using the corresponding assigned subkey~~which subkeys are assigned to one data bit each (Mp[n-1]-Mp[0]) of the permuted data word (Mp), and which each indicate the data bit (M[n-1]-M[0]) of the data word to be permuted (M), which data bit is to be mapped to this data bit (Mp[n-1]-Mp[0]), wherein each subkey (P[n-1]-P[0]) comprises a number of key bits (P[n-1,m-1]-P[n-1,0], P[k,m-1]-P[k,0], P[0,m-1]-P[0,0]).~~

5. (Currently Amended) The Mmethod of according to Cclaim 4, where in which the step of mapping comprises: ~~of a data bit (M[n-1]-M[0]) of the data word to be permuted (M) to a data bit (Mp[k]) of the permuted data word is effected incrementally using a subkey (P[k]) by the following steps:~~

a) selecting a first group of the data bits of the data word to be permuted ( $M_p$ ) as determined by a first one of the plurality of key bits ( $P[k,0]$ ) of the corresponding assigned subkey ( $P[k]$ );

b) selecting a second group of the data bits of the data word from the first group of the data bits obtained by the previous selection as determined by a second one of the plurality of key bits ( $P[k,1]$ ) of the corresponding assigned subkey ( $P[k]$ ); and

c) repeating step b), each time using an additional one of the plurality of key bits of the corresponding assigned subkey ( $P[k,2]...P[k,m-1]$ ) until there exists selected group comprises only one remaining more data bit of the data word, where the one remaining data bit which corresponds to the data bit ( $M_p[k]$ ) of the data word mapped to of the corresponding data bit of the permuted data word ( $M_p$ ).

6. (Currently Amended) The Mmethod of according to Claim 5, where in which the number of data bits contained in the second a group of the data bits of the data word is reduced from one step to the next by a factor of 2two from the number of data bits in the first group of the data bits of the data word, and where the number of data bits in each group of the data bits of the data word in each iteration of step c is reduced by a factor of two.

7. (Currently Amended) The Mmethod according to one of the foregoing claims 2, where in which the first substitution key ( $S$ ) has a number includes a plurality of key bits ( $S[n-1]...S[0]$ ) corresponding to the number of data bits of the permuted of the data word to be substituted ( $M_p$ ), where the step of substituting each data bit of the permuted data word using a substitution key further comprises the step of mapping wherein each data bit of the permuted

data word to be substituted ( $M_p$ ) is mapped unchanged or inverted to a data bit ( $M'[n-1]...M'[0]$ ) of the substituted data word ( $M'$ ) in one of an unchanged form and an inverted form as determined by the corresponding one of these key bits ( $S[n-1]...S[0]$ ).

8. (Currently Amended) The Mmethod according to one of the foregoing claims 3, in which the permutation key (P) and the substitution key (S) are regenerated before a rewriting to the memory after a deletion where the substitution key includes a plurality of key bits corresponding to the number of data bits of the data word, where the step of substituting each data bit of the data word using a substitution key further comprises the step of mapping each data bit of the data word to a data bit of the substituted data word in one of an unchanged form and an inverted form as determined by the corresponding one of the key bits.

9. (Currently Amended) The Mmethod according to one of the foregoing claims 1, further comprising the step of which in order to generating the a permutation key by the (P) comprises the following steps:

a) randomly generating a sub-permutation-key and assigning the generated sub-permutation-key to a data bit position of the permuted data word;

b) checking whether the generated sub-permutation-key has already been assigned to a generated for another data bit position of the permuted data word, and retaining the generated sub-permutation-key as the assigned sub-permutation-key if the generated sub-permutation key it has not yet been assigned to a data bit of the permuted data word generated, and rejecting the generated sub-permutation-key if it has already been generated; and

c) implementing steps a) and b) until a sub-permutation-key is ~~assigned~~<sup>generated</sup> to for each data bit position of the permuted data word ( $M_p$ ).

10. (Currently Amended) The Mmethod of claim 1, further comprising the step of decrypting the stored permuted data word using a second permutation key matched to the permutation key used to generate the permuted data word~~according to one of the foregoing claims, in which a data word ( $M'$ ), generated from a data word ( $M$ ) using the first key after being read out from the memory, is permuted in order to generate the data word using a second permutation key ( $P'$ ) which is matched to the first permutation key ( $P$ ).~~

11. (Currently Amended) A Ddevice that to encrypts and/ decrypts a data word ( $M$ ) comprising having a predetermined number of data bits ( $M[n-1], M[k], M[0]$ ), the which device having has a permutation unit (14) with comprisingthe following features:

—a plurality of data inputs that receive to supply the data bits ( $M[n-1], M[k], M[0]$ ) of the data word to be permuted ( $M$ ); and

—outputs to supply the data bits ( $M_p[n-1], M_p[k], M_p[0]$ ) of a permuted data word ( $M_p$ ) of the predetermined length ( $n$ );

—permutation key inputs to supply a permutation key ( $P$ ) which comprises a number ( $n$ ) of subkeys ( $P[n-1]... P[0]$ ) corresponding to the number of data bits;

—a number plurality of selection units ( $14\_n-1, 14\_k, 14\_0$ ) corresponding to the number of data bits of the data word, whereto which selection units each one of the selection units is responsive to a one-subkey portion of a permutation keyeach is assigned, where each one of the selection units and which provides one data bit each ( $M_p[n-1], M_p[k], M_p[0]$ ) of a the permuted

data word ( $M_p$ ) ~~from the corresponding data bit of the data word as determined by the~~  
~~corresponding one each of the subkeys ( $P[n-1] \dots P[0]$ ) from the data bits of the data word to be~~  
 permuted ( $M$ ).

12. (Currently Amended) ~~The Device of according to Claim 11, where in which each of the~~  
 selection units ( $14_k$ ) ~~comprises has~~ a number of consecutively arranged selection stages  
 ( $141_{n-1}, 141_k, 141_0$ ) ~~corresponding to a~~ the number of permutation key bits ~~of the~~  
~~corresponding subkey for that selection unit, wherein a first selection stage ( $141_0$ ) is designed,~~  
 is responsive to a first one of the permutation key bits ~~as determined by a first key bit ( $P[k,0]$ ), to~~  
 select and provide a first group of data bits ~~offrom the data word to be permuted ( $M$ ), and~~  
 wherein subsequent ~~ones of the~~ selection stages ( $141_1, 141_2, 141_{m-1}$ ) ~~are designed, in each~~  
 case ~~as determined by~~ are each responsive to subsequent ones of the permutation key bits ~~a key~~  
 bit ( $P[k,1], P[k,2], P[k,m-1]$ ), ~~to select a subgroup of the data bits from a the group of data bits of~~  
 the data word provided by the respective previous selection stage.

13. (Currently Amended) ~~The Device of according to Claims 11 or 13, further comprising a~~  
~~in which~~ a substitution unit is connected ~~before or after the permutation unit ( $14$ ), that~~ which  
 substitution unit substitutes ~~each~~ data bits ( $M_p[n-1], M_p[k], M_p[0]$ ) ~~of the permuted a data~~  
 word ~~to be substituted ( $M_p$ ) as determined by in response to a substitution key ( $S$ ).~~

14. (New) The device of claim 11, further comprising a substitution unit connected before  
 the permutation unit, that substitutes each data bit of the data word in response to a substitution  
 key.

15. (New) A method of storing encrypted data in a memory, comprising the steps of:

encrypting a data word by permutating each data bit of the data word using a permutation key to generate a permuted data word;

substituting each data bit of the permuted data word using a substitution key to generate a substitute data word; and

storing the substitute data word in the memory.

16. (New) The method of claim 15, where the permutation key includes a plurality of subkeys corresponding to the number of the data bits of the unencrypted data word, and where each one of the subkeys includes a plurality of key bits, where the step of permutating each data bit further comprises the steps of:

assigning each one of the subkeys to a corresponding one of the data bits of the permuted data word; and

mapping each data bit of the data word to a corresponding one of the data bits of the permuted data word using the corresponding assigned subkey.

17. (New) The method of claim 16, where the step of mapping comprises the following steps:

a) selecting a first group of the data bits of the data word as determined by a first one of the plurality of key bits of the corresponding assigned subkey;



b) selecting a second group of the data bits of the data word from the first group of the data bits as determined by a second one of the plurality of key bits of the corresponding assigned subkey; and

c) repeating step b), each time using an additional one of the plurality of key bits of the corresponding assigned subkey until there exists one remaining data bit of the data word, where the one remaining data bit corresponds to the data bit of the data word mapped to the corresponding data bit of the permuted data word.

18. (New) A method of storing encrypted data in a memory, comprising the steps of:

substituting each data bit of an unencrypted data word using a substitution key to generate a substitute data word; and

permutating each data bit of the substitute data word using a permutation key to generate a permuted data word;

storing the permuted data word in the memory.

19. (New) The method of claim 18, where the permutation key includes a plurality of subkeys corresponding to the number of the data bits of the substitute data word, and where each one of the subkeys includes a plurality of key bits, where the step of permutating each data bit further comprises the steps of:

assigning each one of the subkeys to a corresponding one of the data bits of the substitute data word; and

mapping each data bit of the substitute data word to a corresponding one of the data bits of the permuted data word using the corresponding assigned subkey.

20. (New) The method of claim 19, where the step of mapping comprises the following steps:

a) selecting a first group of the data bits of the substitute data word as determined by a first one of the plurality of key bits of the corresponding assigned subkey;

b) selecting a second group of the data bits of the substitute data word from the first group of the data bits as determined by a second one of the plurality of key bits of the corresponding assigned subkey; and

c) repeating step b), each time using an additional one of the plurality of key bits of the corresponding assigned subkey until there exists one remaining data bit of the substitute data word, where the one remaining data bit corresponds to the data bit of the substitute data word mapped to the corresponding data bit of the permuted data word.